CSE 231—Advanced Operating Systems "Disco"

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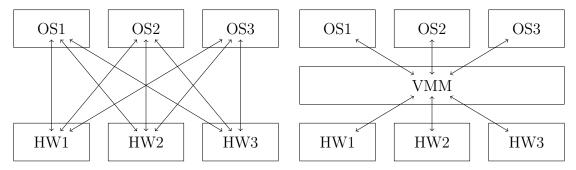
Background. Virtual machines were introduced in the 60s. In the 70s, Popek and Goldberg introduced a distinction between so called type-1 and type-2 hypervisors [5], which remains widely used today. Type-1 hypervisors, sometimes called 'bare-metal', execute directly on top of the hardware and guests are deployed above (see Figure 1a). Prevalent Type-1 hypervisors include VMWare ESXi, and Xen. In contrast, Type-2 hypervisors execute on top of an existing operating system and are thus essentially equivalent to a process in an OS (see Figure 1b). Example Type-2 hypervisors include VirtualBox, VMWare workstation, and QEMU. Finally, some systems do not neatly fit into either category. For example, KVM [] has both Type-1 and Type-2 characteristics: it is a kernel module that essentially makes Linux a Type-1 hypervisor, but deploys guest OSes on top of Linux. The community continues to rigorously research, Type-1, Type-2, and other uncategorized hypervisors,



(a) Type-1 Hypervisor. A Type-1 hypervisor (b) Type-2 Hypervisor. A Type-2 hypervisor (VMM) executing alongside a single application (G1 and G2).
(App1) and running a single guest Operating system (G1).

Summary. Disco investigates the use of Virtual Machine Managers (hypervisors) enabling non-uniform memory access (NUMA) multiprocessors. Their high-level vision is that hypervisors present a solution to "computer vendors attempting to provide system software for innovative hardware", with shared-memory NUMA multiprocessors as the particular innovative hardware that they target. The authors propose a number of key solutions to support a fully virtualized platform, in which a guest experiences no difference between running on raw hardware compared to running on the hypervisor. Disco creates virtualization layers for most devices and computer components, including memory, IO devices, and networking devices. Disco achieves scalability by deploying existing OSes side-by-side and using existing distributed systems protocols for communication (e.g., NFS provides a shared file system across guests). The results are strong, with virtualization coming with relatively meager costs (3%-16%) depending upon workload)

Questions/comments. At first glance, the paper's main motivation seems stellar. A layer of indirection that would allow existing systems to seamlessly adopt innovative hardware is clearly a huge win, because it turns an O(m * n) problem into an O(m + n) problem ¹(see below).



(a) **No Abstraction.** Each Operating System (b) **VMM Abstraction.** Each OS and archiwould need to support each architecture in order tecture only needs to use or support the virtual adopt innovation. machine interface.

The problem is that unlocking the performance potential of innovative hardware will almost invariably require operating system and application customization. For example, developers have updated the design of nearly all major operating systems in light of NUMA and multicore systems. Moreover, the magnitude of hardware innovations that can be effectively masked is necessarily minimal. For example, performance dictates that VMs execute directly on the hardware (i.e., there is no emulation layer), but that requires that all hardware innovation be ISA agnostic.

Nevertheless, virtual machines remain an incredibly important primitive and construct in operating systems today. In class, we identified the following reasons why developers and operators love using virtual machines:

- 1. Fault Domains. Virtual machines separate a fault in one operating system from another. This is especially useful when hacking a custom system.
- 2. **Standardization.** Virtual machines provide standardization in software deployment for both developer and operations teams.
- 3. Security. Virtual Machines, provide separate security domains, at least in principle [].
- 4. Migration and Checkpointing. Virtual Machines are a strong abstraction upon which to implement checkpointing and migration across deployed nodes.
- 5. Fault Tolerance. Virtual Machines are a strong abstraction upon which to build system-level fault tolerance[1].

¹This terminology is from Sky Computing [6]

6. **Record & Replay** Virtual Machines are a strong abstraction upon which to implement deterministic record and replay, which further enables security forensics [2], debugging [4], intrusion detection [3], etc.

References

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