Bugs as Deviant Behavior: A General Approach to Inferring Errors in Systems Code

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Background and motivation

- Different methods used to find errors in a system
 - Testing and manual inspection
 - Type systems
 - Formal verification
 - High-level compilation
 - Dynamic analysis



image reference

Motivation

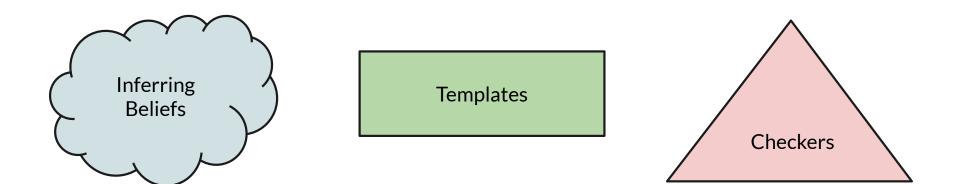
- Often difficult to derive the exact correctness rules for a system
- How can we still design a checker without prior knowledge about the system ?

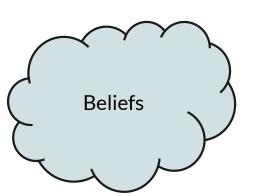


image reference

Design Principles

- Requires no knowledge about system correctness rules
- Infer programmer's beliefs from source code
 - "if two beliefs contradict, we know that one is an error without knowing what the correct belief is."
- If there is a contradiction, then there is *at least one* statement which is wrong.





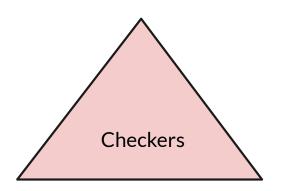
- MUST beliefs, directly implied by the code
- Any contradiction means there is an error in the code



- MAY beliefs, suggested beliefs, could be a coincidence
- Not all contradictions are errors
- Need to separate out noise from errors

Templates

- Outline for a rule
- Example, <a> must be paired with
 - <a> and positions are slots
 - Filled with elements from source code
 - Slot instances, example "lock" and "unlock" function calls.



- General method for finding bugs
- Internal consistency checkers used with MUST beliefs
- Statistical analysis checkers used with MAY beliefs

Framework for internal consistency checkers

- The rule template T
- The valid slot instances for T
- The code actions that imply beliefs
- The rules for how beliefs combine, including the rules for contradictions
- The rules for belief propagation

Reference : from paper

Internal consistency checkers

- MUST beliefs inference
 - Direct observation
 - Implied pre and post conditions
- More beliefs found, more applicable the checker
- Ranking results not necessary because a single contradiction results in an error

Framework for statistical analysis checkers

- It applies the check to all potential slot instance combinations, it assumes that all combinations are MUST beliefs.
- It indicates how often a specific slot instance combination was checked and how often it failed the check (errors).
- It is augmented with a function, *rank*, that uses the count information above to rank the errors from all slot combinations from most to least plausible.

```
// Lock
 1: lock 1;
                    // Variables potentially
 2: int a, b;
                    // protected by 1
 3: void foo() {
                    // Enter critical section
 4:
        lock(1);
        a = a + b; // MAY: a,b protected by 1
 Б:
 6:
        unlock(1); // Exit critical section
 7:
        b = b + 1; // MUST: b not protected by 1
 8: }
 9: void bar() {
10:
        lock(1):
11:
        a = a + 1; // MAY: a protected by 1
12:
        unlock(1):
13: }
14: void baz() {
15:
        a = a + 1; // MAY: a protected by 1
16:
        unlock(1);
17:
        b = b - 1; // MUST: b not protected by 1
18:
        a = a / 5; // MUST: a not protected by 1
19: }
```

Figure 1: A contrived, useful-only-for-illustration example of locks and variables

Statistical analysis checkers

• z statistics for proportions used for sorting between noise and errors

$$z(n,e) = (e/n - p_0)/\sqrt{(p_0 * (1 - p_0)/n)}$$

n = number of check
e = number of successful checks
p₀ = probability of the examples

• Latent specifications to prune the search space

Performance Evaluation

- Analyses written using *Metal* high- level state machine (SM) language for writing system- specific compiler extensions
- Tested on Linux and OpenBSD
 - $\circ \quad \ \ Linux \, 2.4.1 \, and \, 2.4.7$
 - OpenBSD 2.8
- Checkers implemented and tested
 - Internal Null Consistency
 - Security Checker
 - Failure Checker
 - Temporal rules derivation

Performance Evaluation

Internal Null Consistency : Finds pointer errors, flags three types of contradictory or redundant beliefs

Checker	Bug	False
check-then-use	79	26
use-then-check	102	4
redundant-checks	24	10

Security: checks for kernel safe pointers, and "tainted" pointers, raise error if a pointer is both

OS	Errors	False	Applied
OpenBSD 2.8	18	3	1645
Linux 2.4.1	12 (3)	16 (1)	4905
Linux 2.3.99	5	n/a	n/a

What's the overhead associated with so many applications of the checker ?!

Performance Evaluation

Failure checker : Find routines that aren't checked for failures

• Found some unexpected, error - not detected before !! IS_ERR consistency checking

Violation of temporal rules : checking to make sure sequence of actions is followed. One case is making sure, freed memory is not used.

- Made use of latent specifications to prune for applicable function pairs
- Hierarchical ranking for reducing the number of false positives.

Takeaways

- Hundreds of errors discovered in real systems, resulting in kernel patches !!
- Some unexpected, serious bugs discovered too!!
- Fairly higher number of false positives reported

Conclusion

- Automatic inference of bugs without system knowledge
- Presents two checker frameworks that implement this
- Easily re-targetable to new systems and fixed overheads*
- Future works on complete automation using machine learning approaches

A very interesting work, with promising performance and future directions, which could address the issues of existing need for added manual inspection and analysis overheads.

Questions and Discussion

- 1) In a system that is designed as a checker, as in this paper, how would we model and account for "completeness", given they can find bugs but can't guarantee the absence of bugs.
- 2) What if a "belief" doesn't fit a template? How common would these cases be and how scalable/ adaptable is this method in such cases how expensive is it to come up with new templates, or would we have to then come up with other tools in order to analyze such beliefs?
- 3) They mention about augmenting static analysis with dynamic monitoring, which seems very promising. What could be some advantages ? Is this something used today ?